編號: 206

國立成功大學 109 學年度碩士班招生考試試題

系 所:電機資訊學院-資訊聯招

考試科目:程式設計

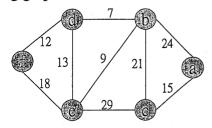
考試日期:0210,節次:2

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※ 考生請注意:本試題不可使用計算機。 請於答案卷(卡)作答,於本試題紙上作答者,不予計分。

Part I. 資料結構 (50%)

- 1. Answer **True** or **False** for the following statements. Give correct answers for **False** statements. (30%, 6% for each question)
 - (1) In static hashing, the worst-case number of comparisons needed for a successful search is O(n) for open addressing. The number could be reduced to $O(\log n)$ by using balanced search tree.
 - (2) Let d_i be the degree of vertex i in a graph G with |V| = n and |E| = e, then $e = \sum_{i=0}^{n-1} d_i$.
 - (3) The path from vertex u to vertex v on a minimal cost spanning tree of an undirected graph G is also a shortest path from u to v.
 - (4) Let G be a graph with e edges and v vertices. If G is represented by adjacency lists, DFS requires $O(ev^2)$ time.
 - (5) If an AOV network represents a feasible project, its topological order is not unique.
- 2. A Bloom Filter consists of m bits of memory and h uniform and independent hash functions f_1 , f_2 , ..., f_h . Each f_i hashes a key k to an integer in the range [1, m]. Initially all m filter bits are zero, and the data set is empty. When key k is added to the data set, bits $f_1(k)$, $f_2(k)$, ..., $f_h(k)$ of the filter are set to 1. When a query "Is key k is in the data set?" is made, bits $f_1(k)$, $f_2(k)$, ..., $f_h(k)$ are examined. The query answer is "maybe" if all these bits are 1. Otherwise, the answer is "no." A filter error occurs whenever the answer is "maybe" and the key is not in the data set. Assume that key k is an integer in the range [1, n] and k updates are made. Compute the probability of filter error for an arbitrary query after the k-th update. (10%)
- 3. Consider the following graph:



Compute minimum cost and construct minimum spanning tree. (10%)

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Part II. 演算法 (50%)

- 1. Is $5^{n+1} = O(5^n)$? Is $6^{2n} = O(6^n)$? (10%)
- 2. Give asymptotic upper and lower bound for $T(n) = \sqrt{n}T(\sqrt{n}) + n$. Assume that T(n) is constant for sufficiently small n. Make your bounds as tight as possible. (10%)
- 3. Show how the time complexity of quicksort can be made to run in $O(n \lg n)$ time. (10%)
- 4. Determine an Longest Common Subsequence of <1,0,0,1,0,1,0,1 and <0,1,0,1,1,0,1,1,0>. (10%)
- 5. If possible, use the master method to solve $T(n) = 27T(\frac{n}{3}) + \Theta(\frac{n^3}{\lg n})$. (10%)