

※ 考生請注意：本試題不可使用計算機。請於答案卷(卡)作答，於本試題紙上作答者，不予計分。

1. In each of the following question, please specify if the statement is true or false.
 - a. [2%] $n^{1.5} = O(n \log n)$
 - b. [2%] A complete binary tree with a height of h can have more nodes than a full binary tree with a height of h .
 - c. [2%] When we use a max heap to implement a priority queue, the time complexity of both the add and delete operations are $O(n)$.
 - d. [4%] $T(n) = T(n - 1) + n$, $T(1) = 1$. Then $T(n) = O(n^3)$
2. In computer science, a priority queue is an abstract data type which is like a regular queue, but each element has a "priority" associated with it. Elements with higher priorities are served before elements with lower priorities.
 - a. [5%] Design a data structure in C++ or Java to represent priority queues where both the data and the priority are integers. Note: 1. Only the class definition is required; 2. The size of the queue should only be constrained by the memory space; 3. Your design should allow efficient enqueue and dequeue operations; 4. The use of STL classes is prohibited.
 - b. [10%] Implement the enqueue and dequeue operations/methods.
3. Answer the following question regarding Red-Black tree
 - a. [2%] Define red-black tree
 - b. [8%] What is the big O performance (in terms of the number of nodes in the tree) for the operations find, insert, and remove for a red-black tree in the best and worst cases?
 - c. [10%] Show the tree that results from inserting the values 2, 1, 4, 5, 9, 3, 6, 7 into an initial empty red-black tree. Show the tree after each insertion.
4. [5%] Given an expression, such as " $\{[2 \times (a + b)] - c\} \times 3(d + e)$ ", implement a function in C++ or Java that checks the proper opening and closing of parenthesis.

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5 · [32%] True or False, and EXPLAIN

Choose T(true) or F(false). If the statement is correct, *briefly state why*. If the statement is wrong, *explain why or give a counterexample*. Answers **WITHOUT** reasons will get **at most 1 point**.

- (a) [4%] (T, F) If we want to randomize the ordering of n integers $A[1], \dots, A[n]$, we can generate another n random numbers $B[1], \dots, B[n]$ with values in $[-n^5, n^5]$, and then sort $B[\cdot]$ in ascending order. We can then use this ordering to reorder $A[1], \dots, A[n]$. Therefore, to randomize n integers takes $\Omega(n \log n)$ time.
- (b) [4%] (T, F) Let $f(x) = \sum_{i=0}^n A[i]x^i$, $g(x) = \sum_{j=0}^m B[j]x^j$, $f(x) + g(x) = \sum_{k=0}^{\max\{m,n\}} C[k]x^k$, and $f(x) \cdot g(x) = \sum_{l=0}^{mn} D[l]x^l$. Calculating $C[]$ and $D[]$, takes $O(\max\{m,n\})$ and $\Theta(mn)$ time, respectively.
- (c) [4%] (T, F) Based on (b), if $A[]$ and $B[]$ store their p and q non-zeros in two ordered doubly linked lists in ascending order. Then, to calculate $C[]$ and $D[]$ takes $O(p+q)$ and $\Theta(pq)$ time, respectively.
- (d) [4%] (T, F) For a binary search tree of n integers, if we insert an integer of new value to the tree, and then delete that integer right away. It is possible the tree structure becomes different from its original form.
- (e) [4%] (T, F) Given a complete undirected graph K_n with n nodes, let $c_{ij} > 0$ represent the length of any edge (i,j) . We can use Dijkstra's algorithm to find a shortest simple path that has to pass all n nodes.
- (f) [4%] (T, F) Given an undirected graph $G=(N,A)$ of $|N|=n$ nodes and $|A|=m$ arcs, let $c_{ij} \geq 0$ represent the length of $(i,j) \in A$. To detect whether there is a zero-length cycle passing node i takes $O(m)$ time.
- (g) [4%] (T, F) Given a Taipei city street map composed by street segments with n nodes and m arcs, where an arc is a street segment with two adjacent nodes. The average degree for a node is $\Omega(n)$.
- (h) [4%] (T, F) To build a min-heap of n values, it takes $\Omega(n \log n)$ time.

6 · [18%] Given a simple directed network $G=(N,A)$ of $n=|N|$ nodes $|A|=m$ arcs, let $c_{ij} \geq 0$, x_{ij} and u_{ij} represent the length, flow, and capacity for $(i,j) \in A$. If n is even and $U = \max_{(i,j) \in A} \{u_{ij}\}$, we want to send a flow of F_k units for $n/2$ pairs $(o_k, d_k) = (k, k+n/2)$, $k = 1, \dots, n/2$. Answer the following questions:

- (a) [6%] Briefly explain how to determine whether F_1 can be successfully sent without violating the arc capacity from the origin node 1 to the destination node $1+n/2$. What is the complexity of your method?
- (b) [6%] Briefly explain how to construct a sub-network $G_k = (N_k, A_k)$ which only contains required nodes and arcs on all possible shortest paths (i.e., the minimum number of arcs) connecting from an origin node o_k to a destination node d_k . What is the complexity of your method?
- (c) [6%] Suppose the complexity of your method in (a) is $S(n,m,U)$ time. To send all the flows of F_k units for each $k = 1, \dots, n/2$ at the same time, briefly explain how to determine whether this is possible or not (if we cannot send F_k units for some k). Is there a method that can do this in $O(S(n,m,U))$ time?