

國立臺灣海洋大學 103 學年度研究所碩士班招生考試試題

考試科目：基礎計算機科學（含資料結構、演算法）

系所名稱：資訊工程學系碩士班不分組

1. 答案以橫式由左至右書寫。2. 請依題號順序作答。

1. (15 %) A matrix with m rows and n columns is called an $m \times n$ matrix. An $m \times n$ matrix is stored in an array in row-major order. Fill the five blanks in the function `void matrix_multiplication(double *a, double *b, double *c, int m, int n, int k)` used to calculate the matrix product: $C=AxB$, where the sizes of matrices **A**, **B**, and **C** are $m \times n$, $n \times k$, and $m \times k$, respectively, and matrices **A**, **B**, and **C** are stored in arrays a , b , and c , respectively. The element C_{ij} in the i th row and the j th column of matrix **C** can be obtained by the formula: $C_{ij} = \sum_u A_{iu}B_{uj}$.

```
void matrix_multiplication(double *a, double *b, double *c, int m, int n, int k)
{
    int i, j;
    for(i = 0; i < m; ++i) {
        for(j = 0; j < k; ++j) {
            double *aPtr = a + (1), *ePtr = aPtr + n;
            double *bPtr = b + (2);
            double *cPtr = c + (3);
            *cPtr = 0;
            for(; aPtr < ePtr; aPtr += (4), bPtr += (5)) {
                *cPtr += *aPtr * *bPtr;
            }
        }
    }
    return;
}
```

2. (15 %) Write a function `int range_count(struct binarytree *rootPtr, int a, int b)` which accepts a pointer `rootPtr` to the root of a binary search tree T and returns the number of nodes in T having key values between a and b ($a \leq key \leq b$). The node structure of the binary search tree is defined as

```
struct binarytree {
    int key;
    binarytree *left, *right;
};
```

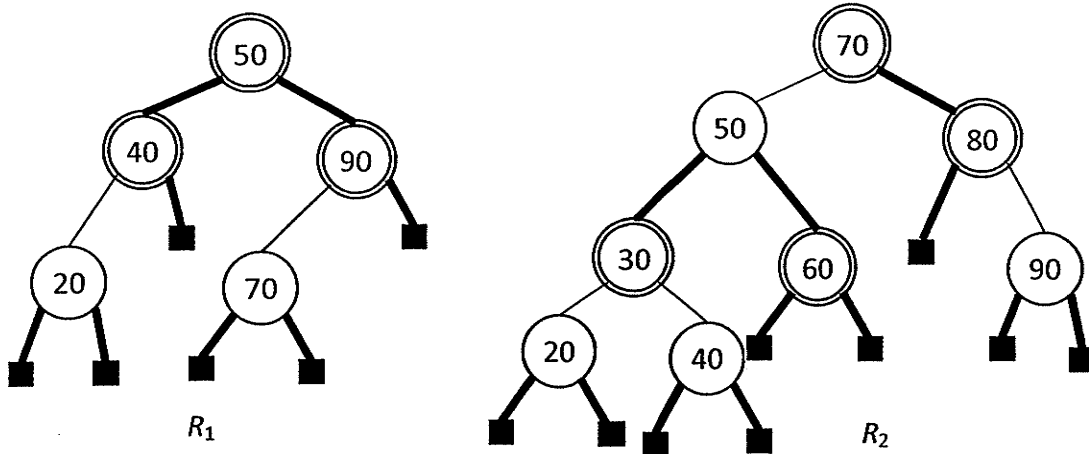
You may call other user-defined functions in the function `range_count`.

3. (10 %) Answer the following questions about the red-black trees R_1 and R_2 .

(a) Show the red-black tree R_1 after insertion of 80.

(b) Show the red-black tree R_2 after deletion of 70.

Notice that the rebalancing rotation and color change are needed to keep the property of the red-black tree.



4. (10 %) There is a huge directed acyclic graph G . The in-degrees of the nodes in G are always less than three, and the out-degrees of the nodes in G are either zero or more than seven. If there is a directed path from node P to node Q , node P is defined as a predecessor of node Q . The adjacency lists and the inverse adjacency lists of G are both available. In order to determine if node P is a predecessor of node Q , we may use a graph search algorithm either to find a directed path from node P to node Q with the adjacency lists of G or to find a directed path from node Q to node P with the inverse adjacency lists of G . Give short answers to the following questions about determining if node P is a predecessor of node Q .

(a) Which of the following graph search algorithms is the better in terms of space complexity? Explain.

the depth-first search
versus
the breadth-first search

(b) If the breadth-first search is used, which of the following search directions is the better in terms of the worst-case time complexity? Explain.

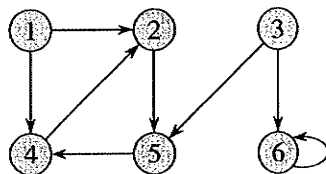
searching from node P with the adjacency lists of G
versus
searching from node Q with the inverse adjacency lists of G

5.(10%) Use a recursion tree to determine a good asymptotic upper bound on the recurrence $T(n) = 2T(n/2) + \theta(n)$.

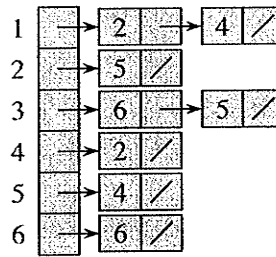
6.(15%) Describe briefly the Quicksort algorithm along with the time complexity. Show how Quicksort can be made to run in $O(n \lg n)$ time in the worst case.

Hint: Selection algorithm.

7.(10%) Illustrate the progresses of BFS and DFS, respectively, starting from vertex 3 on the following graph. Show the state of each phase.



(a)



(b)

8.(15%) The single-source shortest paths problem can be solved by the Bellman-Ford algorithm.

(a) Find the shortest paths starting from vertex 1, going through all other vertices in the following graph by the Bellman-Ford algorithm and show the state of each phase.

(b) Describe briefly the Bellman-Ford algorithm along with the time complexity.

